

What is $\binom{n}{i}$?

$$1) \binom{n}{i} = 0 \quad \text{if } i < 0 \text{ or } i > n$$

$$2) \binom{n}{i} = \frac{n!}{i!(n-i)!}$$

Proof #2

Task: Pick one subset of $[n]$ of cardinality i .

This can be done in " i " steps:

Step 1: Pick 1 of the n elements. (n ways)

Step 2: —||— 1 —||— $n-1$ elements left. ($(n-1)$ ways)

⋮

⋮

⋮

⋮

⋮

Step i : —||— 1 —||— $n-i+1$ elements left. ($(n-i+1)$ ways)

Hence, by the multiplication principle, this task can be done

in $n \cdot (n-1) \cdots (n-i+1) = \frac{n!}{(n-i)!}$ ways.

But! By doing so we obtain $i!$ times each of the i -subsets of n . (By permuting the order of the choices).

Therefore, we divide by $i!$ to get

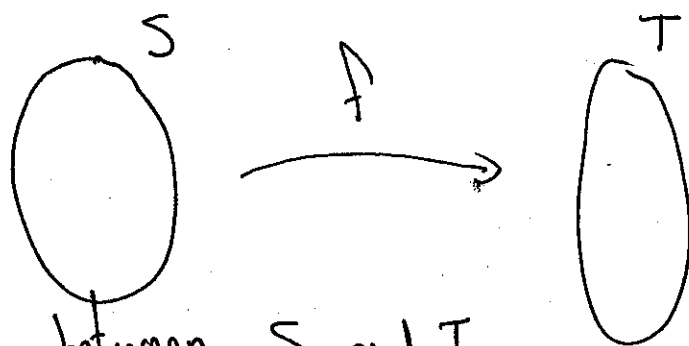
$$\binom{n}{i} = \frac{n!}{i!(n-i)!}$$

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What is a Bijective Proof

D

Two sets:



Find a bijection between S and T.

Examples

$$1) \quad S = \{ A \subseteq [n] \mid |A| = i \}$$
$$T = \{ A \subseteq [n] \mid |A| = n-i \}$$

$$f: S \rightarrow T \quad f \text{ is bijective.}$$
$$A \mapsto [n] \setminus A$$

$$\Rightarrow \binom{n}{i} = \binom{n}{n-i}.$$

$$2) \quad S = \{ A \subseteq [n] : |A| = i \}$$
$$U = \{ A \subseteq [n] : |A| = i, 1 \in A \} \cup \{ A \subseteq [n] : |A| = i, 1 \notin A \}$$

$$g: S \rightarrow U$$
$$A \mapsto A \quad (\text{it is sent to one of the two sets forming } U).$$

g is bijective (check!)

$$\Rightarrow \binom{n}{i} = \binom{n-1}{i-1} + \binom{n-1}{i}.$$

3) Let $n, k, p \in \mathbb{N}$, then

$$\binom{n+k}{n+p} \binom{n+p}{p} = \binom{n+k}{n} \binom{k}{p}$$

Pf - If $k < p$ then both side are 0.

Otherwise, consider

Task T: Select " p " employees from " $n+k$ " applicants, where some subset of the candidates (greater than p) should pass an interview.

Task T can be split into two consecutive tasks as follows:

- 1) - First, select " $n+p$ " applicants for the short list.
- Then, from the short list, select " p " employees.

By the M.P. there are

$$\binom{n+k}{n+p} \binom{n+p}{p} \text{ ways to do this. } \left(\begin{matrix} \text{since} \\ p \leq k \end{matrix} \right)$$

- 2) - First, reject " n " applicants from the " $n+k$ ".
- Then, select " p " employees from the " k " remaining that passed an interview.

By the M.P. there are

$$\binom{n+k}{n} \binom{k}{p} \text{ ways to do this.}$$

We counted in two ways the same task, hence the numbers are equal. \square

(4)

4) For all $n, i \in \mathbb{N}$
$$\sum_{j=0}^n \binom{j}{i} = \binom{n+1}{i+1}$$

Pf Task: Choose a $(i+1)$ -element subset of a $(n+1)$ -element set.

1) The RHS counts this directly $\rightarrow \binom{n+1}{i+1}$ ways.

2) The LHS counts this task in $n+1$ disjoint classes:
The j -th class counts the number of $(i+1)$ -subsets where the largest element is " $j+1$ ".

\hookrightarrow There are $\binom{j}{i}$ elements in the j -th class.
The equality follows from the addition principle. \star

5) (Vandermonde's Sum)

For all $n, m, k \in \mathbb{N}$, we have

$$\binom{n+m}{k} = \sum_{i=0}^k \binom{n}{i} \cdot \binom{m}{k-i}$$

Pf Task: Select a committee of k individuals among " n " bachelor students and " m " master students.

1) This is counted directly by $\binom{n+m}{k}$ (LHS).

The RHS also counts this in $(k+1)$ disjoint sets.

The i -th class counts the committees with " i " bachelor students. There are $\binom{n}{i} \binom{m}{k-i}$ ways to do this by the M.P. The equality holds by the A.P. \square

The Binomial Theorem

Let $n \in \mathbb{N}$. Then
[and $a, b \in \mathbb{R}$]
say

$$(a+b)^n = \sum_{i=0}^n \binom{n}{i} a^i b^{n-i}$$

PF Monomials in the expansion of $(a+b)^n$ are of the form $a^i b^{n-i}$ for some $i \in \{0, 1, \dots, n\}$.

To get $a^i b^{n-i}$, we need to choose " i " brackets where the a 's are coming from; the other terms will give " $n-i$ " b 's.

There are $\binom{n}{i}$ ways to make this choice giving the formula. \square

Applications

1) Show that $\sum_{i=0}^n \binom{n}{i}^2 = \binom{2n}{n}$

PF Consider $(1+x)^n (1+x)^n = (1+x)^{2n}$

By the Binomial Thm:

$$(1+x)^n = \sum_{i=0}^n \binom{n}{i} x^i$$

(6)

The coefficient of x^n on the LHS is obtained by summing the pairings " x^i " with " x^{n-i} " that have coefficients $\binom{n}{i}$ and $\binom{n}{n-i}$ respectively.

Since $\binom{n}{n-i} = \binom{n}{i}$ we get that the coeff of x^n is $\sum_{i=0}^n \binom{n}{i}^2$.

On the other hand, the coefficient of " x " in $(1+x)^{2n}$ is $\binom{2n}{n}$ by the B.T. \star

2) Let $n \in \mathbb{N} \setminus \{0\}$, then

$$\sum_{i=0}^n \binom{n}{i} (-1)^i = 0$$

Prf Let $a=1$ and $b=-1$ in the Binomial Theorem \star

3) What is the coefficient of x^4 in $(3x + 5y + 7)^6$?

Solution: Write $(\underbrace{3x}_a + \underbrace{(5y+7)}_b)^6 = (a+b)^6 \stackrel{\text{B.T.}}{=} \sum_{i=0}^6 \binom{6}{i} a^i b^{6-i}$

The monomial x^4 only appears when $i=4$ in the above sum. Hence the coefficient is

$$\binom{6}{4} a^4 b^2 = \binom{6}{4} (3x)^4 (5y+7)^2 = \binom{6}{4} 3^4 (5y+7)^2 x^4$$

Solution.

4) For all $n \geq 0$,

$$n \cdot 2^{n-1} = \sum_{k=0}^n k \binom{n}{k}$$

Pf From the B.T.,

$$(1+x)^n = \sum_{k=0}^n \binom{n}{k} x^k$$

Differentiating both sides w.r.t. x ,

$$n(1+x)^{n-1} = \sum_{k=0}^n k \binom{n}{k} x^{k-1}$$

Then substitute " x " by "1" to get the desired equality. \square

The Inclusion-Exclusion Principle

Let P_1, P_2, \dots, P_r be properties that elements of a set S can have. Define $A_i := \{a \in S : a \text{ has property } P_i\}$.

We have

$$\left| \bigcup_{i=1}^r A_i \right| = \sum_{i=1}^r |A_i| - \sum_{1 \leq i < j \leq r} |A_i \cap A_j| + \sum_{1 \leq i < j < k \leq r} |A_i \cap A_j \cap A_k|$$

$$\pm \dots + (-1)^{r-1} |A_1 \cap A_2 \cap \dots \cap A_r|.$$

$$= \sum_{\emptyset \neq J \subseteq [r]} (-1)^{|J|-1} \left| \bigcap_{j \in J} A_j \right|.$$

By de Morgan's laws:

$$\begin{aligned}
 \left| \bigcap_{i=1}^r \overline{A_i} \right| &= \left| S - \bigcup_{i=1}^r A_i \right| = |S| - \sum_{i=1}^r |A_i| + \sum_{1 \leq i < j \leq r} |A_i \cap A_j| \\
 &\quad \downarrow \text{def} \\
 &\quad S \setminus A_i \qquad \qquad \qquad -/+ \dots + (-1)^r |A_1 \cap \dots \cap A_r|.
 \end{aligned}$$

Warning: the exponent "r" here is correct. Compare it with the other formula.

- Prf Take $s \in S$.
- If s is not in $\bigcup_{i=1}^r A_i$ then it is not counted in any part of the RHS. (and not on the LHS either...)
 - Else, let $I \subseteq [r]$ be the set of indices i s.t. $s \in A_i$.
 - In the 1st sum, "s" is in $|I|$ summand ($|I| \geq 1$).
 - In the 2nd sum, "s" is in $\binom{|I|}{2}$ —||—
 - ⋮
 - In the $|I|$ th sum, "s" is in $\binom{|I|}{|I|}$ —||—

Hence "s" is counted a total of

$$\binom{|I|}{1} - \binom{|I|}{2} + \binom{|I|}{3} + \dots + (-1)^{|I|-1} \binom{|I|}{|I|} \text{ times.}$$

We can write, after setting $|I| =: k$

$$\begin{aligned}
 -(k-1) &= 1 - \binom{k}{1} + \binom{k}{2} - \binom{k}{3} + \dots + (-1)^{k-1} \binom{k}{k} \\
 &= \sum_{i=0}^k \binom{k}{i} (-1)^{k-i} = (1-1)^k = 0. \quad (\forall k \in \{1, \dots, r\}) \\
 &\Rightarrow k=1 \quad \square.
 \end{aligned}$$

Example

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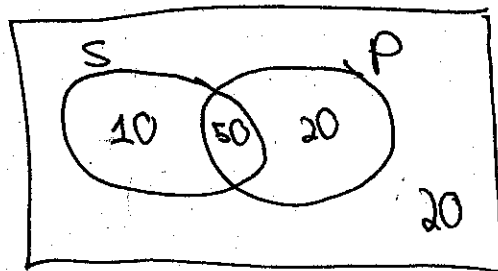
At the last exam

- 60 students studied (Set S)
- 70 " passed (Set P)
- 50 " did both (Set $P \cap S$)
- 20 " did not do either (Set $(P \cup S)^c$)

How many students were at the exam?

Answer:

The class "C"



$$|C| = |(P \cup S)^c| + |P \cup S| = 20 + (60 + 70 - 50) = 100$$

Example: (Dérangements)

A new iPhone virus (or feature?) permutes the phone numbers in the address book. [Assume each entry has one number only] ↳ in the address book contacts

If an iPhone gets the virus and has n contacts, in how many ways can the virus permute the numbers so that no contact gets its own number?

For $n=0$, $D(0) = 1$ $n=2$, $D(2) = 1$

$n=1$, $D(1) = 0$

n=3

Person A

Person B

Person C

B

C

A

C

A

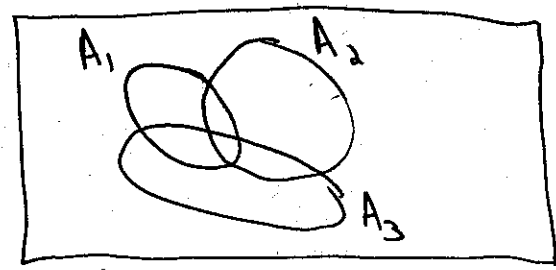
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↳ Method: Go through each permutation and check.
for each n...

... This is long and inefficient.

↳ Use Inclusion-Exclusion principle!

All permutations



Let $S := \{ \text{permutations of } [n] \}$

$A_i := \{ \text{permutations of } [n] \text{ that fix } "i" \}$

~~Let~~ $D := \{ \text{dérangements of } [n] \} = \{ \text{permutations of } [n] \text{ without fixed points} \}$

$$\text{Then } D = \bigcap_{i=1}^n \overline{A_i} = S - \bigcup_{i=1}^n A_i$$

By the Incl.-Excl. principle

$$|D| = |S| - \sum_{i=1}^n |A_i| + \sum_{1 \leq i < j \leq n} |A_i \cap A_j| - \dots + (-1)^n |A_1 \cap \dots \cap A_n|$$

$|A_i| = (n-1)!$ because " i " is fixed and the rest is permuted. (12)

$|A_i \cap A_j| = (n-2)!$ because 2 elements are fixed — // —

Therefore

$$|D| = n! - \sum_{i=1}^n (n-1)! + \sum_{1 \leq i < j \leq n} (n-2)! - / + \dots + (-1)^n \cdot 1$$

$$= n! - n! + \binom{n}{2} (n-2)! - / + \dots + (-1)^n \cdot 1$$

$$= n! - n! + \frac{n!}{2!} - / + \dots + (-1)^n \frac{n!}{n!}$$

$$= n! \left(\frac{1}{0!} - \frac{1}{1!} + \frac{1}{2!} - \frac{1}{3!} + \dots + (-1)^n \frac{1}{n!} \right)$$

$$= n! \sum_{i=0}^n \frac{(-1)^i}{i!}$$

Example: (Euler's totient function)

For $n \in \mathbb{N}$, let $\varphi(n) = \#$ integers m such that $m < n$ and $(m, n) = 1$.

$$\varphi(n) = n \cdot \prod_{\substack{p|n \\ p \text{ prime}}} \left(1 - \frac{1}{p} \right)$$

Pf Let $P = \{p_i \mid p_i \mid n\}$ and $|P| = k$ (12)

and $A_i = \{j : j \leq n \text{ and } p_i \mid j\}$

We want $|[n] \setminus \bigcup_{i=1}^k A_i|$

$$= n - \sum_{i=1}^k |A_i| + \sum_{1 \leq i < j \leq k} |A_i \cap A_j| - \dots + (-1)^k |A_1 \cap \dots \cap A_k|$$

$$= n - \sum_{i=1}^k \frac{n}{p_i} + \sum_{i < j} \frac{n}{p_i p_j} - \dots + (-1)^k \frac{n}{p_1 \dots p_k}$$

$$= n \cdot \left[1 - \sum_{i=1}^k \frac{1}{p_i} + \sum_{i < j} \frac{1}{p_i p_j} - \dots + (-1)^k \frac{1}{p_1 \dots p_k} \right]$$

$$= n \left[\prod_{p \in P} \left(1 - \frac{1}{p} \right) \right]$$

□

The Pigeonhole Principle

(13)

Theorem (Pigeonhole Principle)

No set of the form $[n]$ is equinumerous to a proper subset of itself, where $n \in \mathbb{N}$.

pf

Claim: If $f: [n] \rightarrow [n]$ is injective then it is surjective (thus bijective).

Suppose $[n]$ is in bijection with a proper subset of itself.

That is, there exists $g: [n] \hookrightarrow [n]$
where A is a proper subset of $[n]$
 $\text{Im}(g)$

By the claim g is also surjective.

\hookrightarrow contradiction with the fact that A is a proper subset. \square

Take Home:

"If m objects are distributed into n boxes, and $m > n$, then at least one box receives at least two objects."

Theorem (Sharper Pigeonhole Principle)

Let $f: N \rightarrow R$ be a function with $|N| = n > r = |R|$.

Then, $\exists a \in R$ such that $|f^{-1}(a)| \geq \lfloor \frac{n-1}{r} \rfloor + 1$.

pf If $|f^{-1}(a)| \leq \lfloor \frac{n-1}{r} \rfloor, \forall a \in R$, then

$$n = \sum_{a \in R} |f^{-1}(a)| \leq r \lfloor \frac{n-1}{r} \rfloor < n. \quad \hookrightarrow \square$$

s_i receives label (a_i, b_i) where
 $a_i =$ length of the longest increasing seq. ending at s_i .
 $b_i =$ ———— // ———— decreasing ———— // ————

For $i < j$, then

either $s_i < s_j \Rightarrow a_i < a_j$.

or $s_i > s_j \Rightarrow b_i < b_j$.

\Rightarrow all labels are distinct!

We have $n^2 + 1$ distinct labels hence one of them has to contain an entry $> n$ by the pigeonhole principle. \star

3) Chinese Remainder Theorem

Let m and n be relatively prime positive integers.

The system

$$\begin{cases} x = a \pmod{m} \\ x = b \pmod{n} \end{cases}$$

has a solution for x .

pf | Consider the numbers

$$a, m+a, 2m+a, \dots, (n-1)m+a$$

n numbers that are $\equiv a \pmod{m}$.

Consider then $im+a$ and $jm+a$ with $0 \leq i < j < n$.

Suppose that $im + a = jm + a \pmod n$.

$$\Rightarrow im + a = q_i n + r \pmod n$$

$$jm + a = q_j n + r \pmod n$$

for some q_i, q_j, r .

Subtracting: $(j-i)m = (q_j - q_i)n$.

Since $(m, n) = 1 \Rightarrow n \mid (j-i) \Rightarrow i=j$. \downarrow (with the choice of i, j).

Every number in 0 to $n-1$ appears as remainder mod n , in particular, "b" does. \star

4) Ramsey's Thm (a first version)

In any group of six people there are 3 people that mutually know each other or three people that do not know each other.

Prf Take one person in the group, say "A".
The other 5 either know or don't know "A".

By the sharper pigeonhole principle either

- i) 3 don't know A, $\{ = \text{group "B"}$
- or ii) 3 know A.

Case i) - If 2 among group "B" don't know each other,
3 people mutually don't know each other.

-Otherwise they form 3 people that know each other.
 \downarrow
group B

Case ii) is similar. \star

5) Lossless data compression algorithm

"Lossless data compression algorithm cannot guarantee compression for all input data sets."

That is, "For lossless d.c.a., there is a data set that does not get smaller^{any}."

PN Let A be a lossless d.c.a. and $S = \{f : f \text{ a finite binary file}\}$
be some set of files encoded as binary strings.
 $\hookrightarrow f$ is a sequence of 0's and 1's

Assume: 1) " A " compresses files in S such that

"new length of $f \leq$ length of f "

2) $\exists f \in S$, such that "new length of $f <$ length of f ".

Let $s \in S$ be the shortest file satisfying 2).

Let M be the length of s .

Let N be the compressed length of s , i.e. $N < M$.

Together with s , there are $2^N + 1$ files compressed into 2^N files of length N .

By the pigeonhole principle, there are two files which are compressed to the same output.

\hookrightarrow " A " can not be lossless: how to recover the two files from one compressed?

That means that lossless d.c.a. have to make some files longer!