



# History

- Brlek, Mendès France, Robson & Rubey, *Cantorian Tableaux and Permanents*, L'Enseignement Mathématique (2004)
- Mendès France, *Cantorian Tableaux revisited*, Funct. Approx. Comment. Math. (2007)
- Brlek, Labbé, Mendès France, *Combinatorial variations on Cantor's diagonal*, (submitted), (2011) preprint available on [arxiv.org](http://arxiv.org)

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$s$	$s^{-1}$	$s^{-2}$	$s^{-3}$	$s^{-4}$	$s^{-5}$	$\dots$
0	$a_{11}$	$a_{12}$	$a_{13}$	$a_{14}$	$a_{15}$	$\dots$
0	$a_{21}$	$a_{22}$	$a_{23}$	$a_{24}$	$a_{25}$	$\dots$
0	$a_{31}$	$a_{32}$	$a_{33}$	$a_{34}$	$a_{35}$	$\dots$
0	$a_{41}$	$a_{42}$	$a_{43}$	$a_{44}$	$a_{45}$	$\dots$
0	$a_{51}$	$a_{52}$	$a_{53}$	$a_{54}$	$a_{55}$	$\dots$
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We create the number  $b = b_1 b_2 b_3 b_4 b_5 \dots$  where  $b_i \neq a_{ii}$

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Naturally, we define the permanent of a tableau  $T$

## Definition (Brek et al. (2004))

The *permanent* of a tableau  $T$  is the set of words

$$\text{Perm}(T) = \bigcup_{\pi \in S_n} a_{\pi(1)1} a_{\pi(2)2} \cdots a_{\pi(n)n}$$



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## Definition (Brlek et al. (2004))

A tableau  $T$  is *Cantorian* if no row-words appear in  $\text{Perm}(T)$ , i.e.

$$L \cap \text{Perm}(T) = \emptyset.$$

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Theorem (Brlek et al. (2004))

*Let  $\mathbb{Q} \subseteq L$  be a countable set in  $[0, 1]$  and  $T$  the tableau obtained by the development of the elements in  $L$  in base  $s \geq 2$ . The tableau  $T$  is Cantorian. Meaning :*

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## Theorem (Brlek et al. (2004))

If  $s = 2$ , then we have

$$\text{Perm}(T) = [0, 1] \setminus L.$$

So, if  $L$  contains all algebraic numbers of  $[0, 1]$ ,  $\text{Perm}(T)$  is exactly the set of all transcendental numbers in  $[0, 1]$ .

# Examples - finite tableaux

We note the set of  $n \times n$  tableaux over the alphabet  $A$  containing  $s$  letters by  $\mathcal{T}_n^s$ .



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# Sufficient condition

Proposition (Brek et al. (2004))

*If for every row  $i$ , there exist another row  $i'$  such that  $t_{ij} \neq t_{i'j}$ , for all  $j$ , then the tableau is Cantorian.*

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$\Rightarrow$  The tableau is Cantorian

# Equivalence relation on tableaux

## Remark (Brlek et al. (2004))

*The property of being « Cantorian » is invariant :*

- ▶ *by permuting rows ;*
- ▶ *by permuting columns ;*
- ▶ *given a bijection of the alphabet, replacing the elements of a column by their image under this bijection.*

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# Equivalence relation on tableaux

## Definition

Let  $T', T \in \mathcal{T}_n^s$ . We write

$T' \sim T \iff T'$  can be obtain from  $T$  by a finite sequence  
of invariant transformations.

We will say that  $T'$  is equivalent to  $T$ .

# Parikh composition

Let  $\mathbb{N} = \{0, 1, 2, 3, \dots\}$  and  $\mathbb{N}^* = \{\text{finite words in } \mathbb{N}\}$ .

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## Definition

Let  $w \in A^*$  be a word of length  $n$ . The *Parikh composition*  $\mathfrak{p}_w := \mathfrak{p}(w)$  of  $w$  is a composition of weight  $n$  and of length  $s$  obtain by the function

$$\begin{aligned} \mathfrak{p} : A^* &\rightarrow \mathbb{N}^* \\ w &\mapsto |w|_{\alpha_1} |w|_{\alpha_2} \cdots |w|_{\alpha_s}. \end{aligned}$$

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We extend the function  $\mathfrak{p}$  to tableaux  $\mathcal{T}_n^s, \mathfrak{P} : \mathcal{T}_n^s \rightarrow (\mathbb{N}^*)^n$ .

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(210, 201, 111)



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$(210, 201, 111)$      $(222, 312, 231, 510, 222, 231)$

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# Orders on $\mathbb{N}^*$ and $A^*$

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$$5 \preceq 41 \preceq 32 \preceq 23 \preceq 14 \preceq 311 \preceq 221 \preceq 212 \preceq 131 \preceq 122 \preceq 113$$

Finally, we define a total order  $\blacktriangleleft$  on  $A^*$ , which we call *Parikh composition order* on  $A^*$ .

## Definition

Let  $w, w' \in A^*$ . We write  $w \blacktriangleleft w'$  if and only if

$$\mathfrak{p}_w \prec \mathfrak{p}_{w'} \text{ or } (\mathfrak{p}_w = \mathfrak{p}_{w'} \text{ and } w \leq w').$$

# Total order on tableaux

## Definition

Let  $T, T' \in \mathcal{T}_n^s$ . We extend naturally the order on tableaux

$$T \triangleleft T' \iff c_1 \triangleleft c'_1$$

or  $(c_1 = c'_1, \text{ and } c_2 \triangleleft c'_2)$   
etc.

where  $c_i$  is the  $i$ -th column-word of  $T$  and similarly for  $c'_i$  with  $T'$ .

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## Example

$$\begin{bmatrix} 1 & 1 & 1 \\ 1 & 1 & 1 \\ 1 & 2 & 2 \end{bmatrix} \triangleleft \begin{bmatrix} 1 & 1 & 1 \\ 1 & 1 & 1 \\ 1 & 2 & 3 \end{bmatrix} \triangleleft \begin{bmatrix} 1 & 1 & 1 \\ 1 & 2 & 3 \\ 1 & 1 & 1 \end{bmatrix} \triangleleft \begin{bmatrix} 1 & 1 & 1 \\ 1 & 3 & 2 \\ 1 & 1 & 1 \end{bmatrix} \triangleleft \begin{bmatrix} 2 & 3 & 1 \\ 2 & 2 & 2 \\ 2 & 3 & 1 \end{bmatrix}$$



# Reduced tableaux

## Lemma

Let  $T \in \mathcal{T}_n^s$ . In the class  $[T]$ , there exist a unique minimal representative  $T_{\blacktriangleleft}$ .

## Definition

A tableau  $T' \in [T]$  is *reduced* if its Parikh compositions are equal to the ones of  $T_{\blacktriangleleft}$ .

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## Remark

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$T$ 
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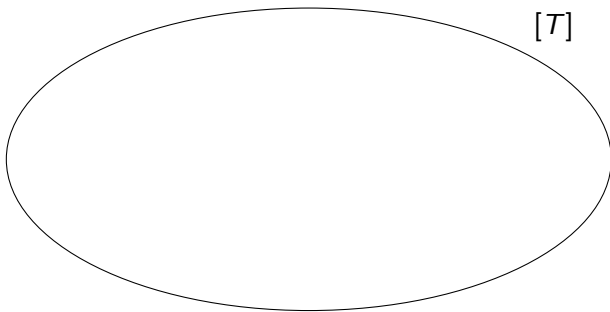
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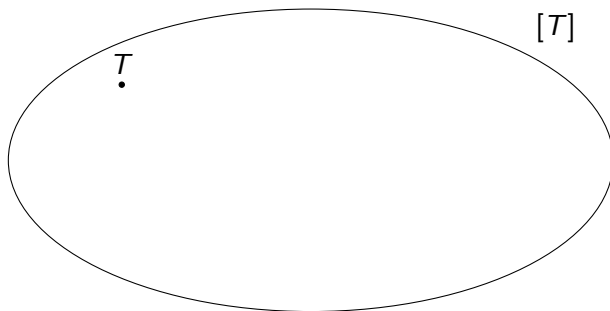




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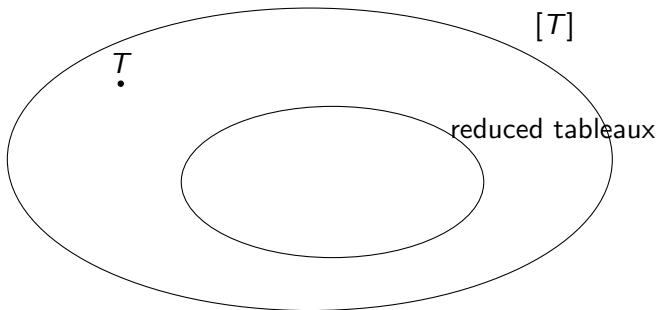
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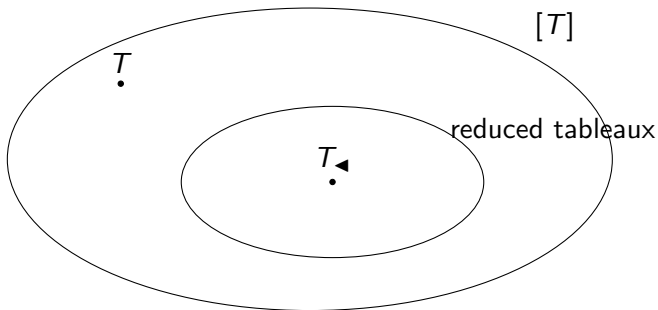
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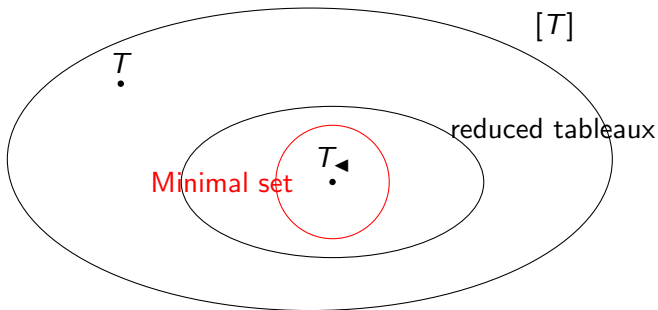
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# Number of Cantorian classes

Number of minimal representatives of dimension  $n$  over an alphabet of  $s$  letters (over the number of classes tested) :

$n \backslash s$	2	3	4	5	6
2	1/1	1/1	1/1	1/1	1/1
3	1/3	5/9	5/9	5/9	5/9
4	6/21	56/171	107/275	107/275	107/275
5	11/165	1873/12574			

# Cantorian minimal representatives

A few minimal representatives

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Dimension  $n = 2$  with  $s \geq 2$  :

$$R_s = \begin{bmatrix} a & a \\ b & b \end{bmatrix}$$
$$|[R_s]| = s^2(s-1)^2$$

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The sufficient condition applies.



# Cantorian minimal representatives

A few minimal representatives

Dimension  $n = 3$ ,  $s = 2$  :

$$R = \begin{bmatrix} a & a & a \\ a & a & a \\ b & b & b \end{bmatrix}$$

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Dimension  $n = 3$ ,  $s = 3$  :

$$\begin{bmatrix} a & a & a \\ a & a & a \\ b & b & b \end{bmatrix}$$

$$|[R_1]| = 648$$

$$\begin{bmatrix} a & a & a \\ a & a & b \\ b & b & c \end{bmatrix}$$

$$|[R_2]| = 1944$$

$$\begin{bmatrix} a & a & a \\ a & b & b \\ b & c & c \end{bmatrix}$$

$$|[R_3]| = 1944$$

$$\begin{bmatrix} a & a & a \\ a & b & b \\ a & c & c \end{bmatrix}$$

$$|[R_4]| = 324$$

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The sufficient condition does not applies !

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## Theorem

Let  $T \in \mathcal{T}_n^s$ . We note the multiplicities of row-words of  $T$  by  $(f_1, f_2, \dots, f_q)$ , where  $q = |L|$ . Similarly,  $(g_1, g_2, \dots, g_r)$  denote the multiplicities of column-words of  $T$ , where  $r = |C|$ .

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The cardinality of  $[T]$  is

$$\frac{(n!)^2}{\left(\prod_{i=1}^r g_j! \prod_{i=1}^q f_i! + \eta\right)} \cdot \prod_{i=1}^n \frac{s!}{(s - \ell^+(p_{c_i}))!},$$

where  $\vartheta = |\mathcal{O}_B(T) \cap \mathcal{O}_\Phi(T)|$  and

$\eta = |\{(\sigma, \tau) \in S_n \times S_n \mid \sigma T \tau = T \text{ and } \sigma T \neq T\}|$ .

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$$|[R_4]| = \frac{\frac{(3!)^2}{((2!1!) \cdot (1!1!1!)) + 0}}{6} \cdot \frac{3!}{(3-1)!} \frac{3!}{(3-3)!} \frac{3!}{(3-3)!} = 324$$

# New enumerative results

Number of Cantorian tableaux of dimension  $n$  over an alphabet of  $s$  letters :

$n \backslash s$	2	3	4	5
2	$1 \cdot 2^2$	$2^2 \cdot 3^2$	$3^2 \cdot 4^2$	$4^2 \cdot 5^2$
3	$3 \cdot 2^3$	$47 \cdot 2^2 \cdot 3^3$	$207 \cdot 3^2 \cdot 4^3$	$579 \cdot 4^2 \cdot 5^3$
4	$109 \cdot 2^4$	$25036 \cdot 2^2 \cdot 3^4$	$803613 \cdot 3^2 \cdot 4^4$	$9419224 \cdot 4^2 \cdot 5^4$
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	Before	After
$C(4, 2)$	3 min.	9 sec.
$C(4, 3)$	$\geq 17\text{h}$	105 sec.
$C(4, 4)$	$\geq 74\text{d}$	25 sec.

Computed using Sage4.5.2  
on a Intel 2.8Ghz machine

# New enumerative results

## Theorem

*The number of Cantorian tableaux  $C(n, s)$  for  $n = 2, 3$  and 4 is given by the following polynomials*

$$C(2, s) = s^2 \cdot (s - 1)^2;$$

$$C(3, s) = s^3 \cdot (s - 1)^2 \cdot (s^4 + 2s^3 - 15s^2 + 16s - 1);$$

$$C(4, s) = s^4 \cdot (s - 1)^2 \cdot (s^{10} + 2s^9 + 3s^8 - 92s^7 - 43s^6 + 1014s^5 - 449s^4 - 5680s^3 + 12045s^2 - 9406s + 2629).$$

# Definition

## Definition (Brlek et al. (2004))

A tableau  $T$  is *bi-Cantorian* if no row-words or column-words appear in  $\text{Perm}(T)$ , i.e.

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## Fact

The property of being « bi-Cantorian » is not invariant under  $\sim$  anymore! ☹



# Brute force enumerative results

$n \setminus s$	2	3	4	5	6	...
2	1 · 2 · 1	2 · 3 · 3	3 · 4 · 7	4 · 5 · 13	5 · 6 · 21	...
3	1 · 2 · 3	2 · 3 · 367	3 · 4 · 6179	4 · 5 · 43065		
4	1 · 2 · 91	2 · 3 · 402873				
5	1 · 2 · 2005					
6						

**TABLE:** Number of bi-Cantorian tableaux of dimension  $n$  over an alphabet of  $s$  letters

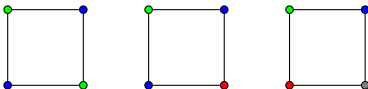
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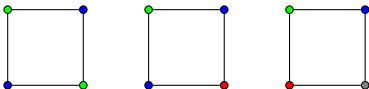


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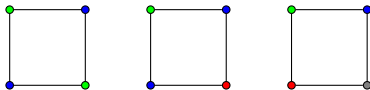


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## Proposition

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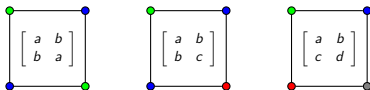
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*Consider an infinite tableau  $T^\infty$  formed by developing in base  $s$  algebraic numbers of  $[0, 1]$ . Are there algebraic columns? If so, how many?*

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Find  $\lim_{n \rightarrow \infty} \frac{B(n,s)}{C(n,s)}$ .

(For  $s = 2$ , the first values are : 0.5, 0.25, 0.104, 0.045)

Merci ! Thank you ! Grazie ! Danke ! Gracias !  
I am founded by :

